

BASIC-ABSTRACT:

NOVELTY - Method consists in identifying all the objects in the program, creating a list of conditions under which objects of the program are to be migrated, and identifying each programmed entity to be accessed from outside of the object. Two proxies are generated for each object, one on the same physical device and the other transferred to the remote computer on object migration. The first proxy contains network linkage and indication to access the programmed entities on the remote computers, and the second contains network linkage and indication to access the programmed entities on the first computer.

Ref #	Hits	Search Query	DBs	Default Operator	Plurals	Time Stamp
L1	2	"6182275".pn.	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 15:43
L2	2	"6,167,383".pn.	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 15:43
L3	1	"6,167,383".pn. and http	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 15:45
L4	0	"6,167,383".pn. and (http with (content field))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 15:45
L5	895	(@ad<"20000410").ad. and (http) and (secure with (box computer machine))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 15:58
L6	430	(@ad<"20000410").ad. and (http) and (secure with (box computer machine)) and (select\$5 same (component part feature))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 15:59
L7	380	(@ad<"20000410").ad. and (http) and (secure with (box computer machine)) and ((select\$5 choos\$5 determin\$5 chos\$5) with (component part feature))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON .	2007/05/08 16:01
L8	380	(@ad<"20000410").ad. and (http) and (secure with (box computer machine)) and ((select\$5 choos\$5 determin\$5 chos\$5) with (component part feature))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:01
L9	Vill 11	(@ad<"20000410").ad. and (http) and (secure with (box computer machine)) and (((select\$5 choos\$5 determin\$5 chos\$5) with (component part feature)) same (configur\$5 assembl\$5))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:02

ProQuest

Return to the USPTO NPL Page | Help

Γ	Bas	sic	1	Advanced	Topics		Publications	7	My Re	search	Interface language
L Da	tah	ase	 2S_S(elected: Multi	ple databases				<u>0 marked</u>	items	
	····	uoc	<u>,,, ,, ,, ,, ,, ,, ,, ,, ,, ,, ,, ,, ,,</u>	orded.	pic databases	•••				•	
R	es	ul	ts								
					PDN(<04/10/2 and HTTP and		00) and (proxy header)	ar	nd (referer	or referre	rer) » Refine Search Set Up Alert 🖂
_				lications			,				
		Ma all	ark	O mar Export	ked items: Er	na	ail / Cite /		Show tex	-	Sort results by: Most recent first
	j	1.	No Av	ne of your e	<u>-business</u> Web Technic	Įu	es. Apr 2000.	Vo	l. 5, Iss. 4;	p. 55 (4	pages)
				∄ <u>Text+</u>	<u>Graphics</u>		🔁 <u>Full</u>	Τe	ext - PDF		^{ta} <u>Abstract</u>
		2.	<u>Ma</u> Je:	nintaining se	cure Web ar	op Di	lications Iting, Manhas	sei	: Mar 20. 2		ol. 11, Iss. 5; p. 93 (4 pages)
		-			<u>Graphics</u>	Γ.			ext - PDF		© Citation
12		3.		e Web is no		ia	ues. Sep 1999		/ol / lee	0·n 18 /	(2 nages)
			LII		Graphics	14			ext - PDF	o, p. 10 (Abstract
		4.	Ka	w safe is the ren Kenworth 144 (9 pages	hy, Nancy A L	.a	ng, Fred Lang	a. V	Windows !	Magazin	ne. Manhasset: Dec 1998. Vol. 9, Iss. 12;
											<u>Citation</u>
	j	5.	An An	evaluation:	Exploring V Capacity Ma	Ve na	b server perfe agement Revi	orr ew	nance . Naples: 0	Oct 1996.	5. Vol. 24, Iss. 10; p. 15 (8 pages)
				Text+	<u>Graphics</u>		🔁 <u>Full</u>	Τe	ext - PDF		[™] Abstract
1-	5 of	5	•	rement to the second	perior en en en en en		The state of the s		eccent i tuet gi interni	errania. V (Principus H	official consistency and form continuous countries of the section of the following of the continuous countries of
W	ant	to	be	notified of n	ew results fo	or	this search?	Se	t Up Alert	×	Results per page: 30 🔽
										,	
В	as	ic	Se	earch			. (Tools:	<u>S</u>	earch Tips	Browse	e Topics 5 Recent Searches
	PI	DN	(<0	4/10/2000) a	nd (proxy an	d	(referer or re	fer	rer) and (s	ecure or	r securi Search Clear
	Da	itab	ase	e: Mülti	ple database	es				386	Select multiple databases
	Da	ite i	ran	ge: All.da	ates		E				
	Lir	nit	resi	ults to: 🔟 Fu	ıll text docum	eı	nts only 🖺		,		
				I Sc	cholarly journa	al	s, including pe	er-	reviewed f	About	<u>ıt</u>

				,		
L10	35	(@ad<"20000410").ad. and ((http) same secure same (box source computer machine)) and (((select\$5 choos\$5 determin\$5 chos\$5) with (component part feature)) same (configur\$5 assembl\$5))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:04
L11	38	(@ad<"20000410").ad. and ((http) same secure same (box terminal source computer machine)) and (((select\$5 choos\$5 determin\$5 chos\$5) with (component part feature)) same (configur\$5 assembl\$5))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:05
L12	17	(@ad<"20000410").ad. and ((http) same access\$5 same secure same (box terminal source computer machine)) and (((select\$5 choos\$5 determin\$5 chos\$5) with (component part feature)) same (configur\$5 assembl\$5))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR ·	ON	2007/05/08 16:07
L13	18	(@ad<"20000410").ad. and ((http) same access\$5 same secure same (box terminal source computer source destination machine)) and (((select\$5 choos\$5 determin\$5 chos\$5) with (component part feature)) same (configur\$5 assembl\$5))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:07
L14	Worm (8)	(@ad<"20000410").ad. and ((http) same access\$5 same(source destination) same secure same (box terminal source computer machine)) and (((select\$5 choos\$5 determin\$5 chos\$5) with (component part feature)) same (configur\$5 assembl\$5))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:33
L15	(2)	(@ad<"20000410").ad. and (http same referrer same proxy)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:35
L16	139	(@ad<"20000410").ad. and (http same header same proxy)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:36
L17	78	(@ad<"20000410").ad. and ((http adj3 header) same proxy)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:36

			-			
L18	23	(@ad<"20000410").ad. and ((http adj3 header) same (proxy and url))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:37
L19	3	(@ad<"20000410").ad. and ((http adj3 header)) and (proxy same (physical logical) same url)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:39
L20	0	(@ad<"20000410").ad. and ((http adj3 header)) and (proxy same (physical logical) same url same source)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:40
L21	1 .	(@ad<"20000410").ad. and ((http adj3 header)) and (proxy same (physical logical) same source same (source URL destination))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:41
L22	50	(@ad<"20000410").ad. and (proxy same (physical logical) same source same (source URL destination))	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:41
L23	14	(@ad<"20000410").ad. and (proxy same (physical logical) same source same (source URL destination)) and http	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:43
L24	14	(@ad<"20000410").ad. and (proxy same (physical logical) same source same (source URL destination)) and http and (known secure)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 17:13
L25	14	(@ad<"20000410").ad. and ((proxy proxies) same (physical logical) same source same (source URL destination)) and http and (known secure)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	ON	2007/05/08 16:44
L26	britis.	(@ad<"19990406").ad. and (custom or customize) and (store near5 front) and (proxy proxies)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/05/08 16:52
L27	2	(@ad<"19990406").ad. and http and (custom or customize) and (store near5 front) and (proxy proxies)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/05/08 16:53

		4				
L28	2	"6,167,383".pn.	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/05/08 16:52
L29	2	(@ad<"19990406").ad. and http and access\$5 and (custom or customize) and (store near5 front) and (proxy proxies)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2007/05/08 16:54
L30	2	(@ad<"19990406").ad. and http and access\$5 and (custom or customize) and (store near5 front) and (proxy proxies)	US-PGPUB; USPAT; EPO; JPO; DERWENT; IBM_TDB	OR	OŅ	2007/05/08 17:13
L31	0	(@ad<"19990406").ad. and http and access\$5 and (custom or customize) and (store near5 front) and (proxy proxies)	EPO; JPO; DERWENT	OR	ON	2007/05/08 17:13
L32	0	(@ad<"20000410").ad. and (proxy same (physical logical) same source same (source URL destination)) and http and (known secure)	EPO; JPO; DERWENT	OR	ON	2007/05/08 17:13
L33	المر	(@ad<"20000410").ad. and (proxy and (physical logical) and source and (source URL destination)) and http and (known secure)	EPO; JPO; DERWENT	OR ·	ON	2007/05/08 17:14
L34	Long (2)	(@ad<"20000410").ad. and (proxy and (physical logical) and (source URL destination))	EPO; JPO; DERWENT	OR	ON	2007/05/08 17:14

DETAILED DESCRIPTION - There is an INDEPENDENT CLAIM for a computer readable code for dynamic distribution of program objects.

USE - Method is for use in moving objects dynamically from computer to computer for load balancing and clustering of objects which communicate frequently to reduce network traffic and improve overall system performance.

ADVANTAGE - Method allows dynamic reconfiguration of programs without programmer intervention, does not require modification of system program source text, and allows the administrator to specify the conditions under which reconfiguration is to occur on a per user basis.

DESCRIPTION OF DRAWING(S) - The figure shows proxy and byte code generation with

object 401

proxy generator 403

local proxy 405

remote proxy 407

byte code modifier 409

serializable object 411

ABSTRACTED-PUB-NO:

US 6324543B

EQUIVALENT-ABSTRACTS:

NOVELTY - Method consists in identifying all the objects in the program, creating a list of conditions under which objects of the program are to be migrated, and identifying each programmed entity to be accessed from outside of the object. Two proxies are generated for each object, one on the same physical device and the other transferred to the remote computer on object migration. The first proxy contains network linkage and indication to access the programmed entities on the remote computers, and the second contains network linkage and indication to access the programmed entities on the first computer.

DETAILED DESCRIPTION - There is an INDEPENDENT CLAIM for a computer readable code for dynamic distribution of program objects.

USE - Method is for use in moving objects dynamically from computer to computer for load balancing and clustering of objects which communicate frequently to reduce network traffic and improve overall system performance.

ADVANTAGE - Method allows dynamic reconfiguration of programs without programmer intervention, does not require modification of system program

source text, and allows the administrator to specify the conditions under which reconfiguration is to occur on a per user basis.

DESCRIPTION OF DRAWING(S) - The figure shows proxy and byte code generation with

object 401

proxy generator 403

local proxy 405

remote proxy 407

byte code modifier 409

serializable object 411

CHOSEN-DRAWING:	Dwg.4/7
TITLE-TERMS:	DYNAMIC DISTRIBUTE PROGRAM OBJECT
DERWENT-CLASS: TO) <u>1</u> ,
EPI-CODES: T01	-F02C; T01-H07C5A; T01-M02A1; T01-S01B;
SECONDARY-ACC-NO:	The state of the s
Non-CPI Secondary Accession N	mbers: N1999-358420

Return to the USPTO NPL Page | Help

Basic Advanced Topics Publications	Interface language
<u>Databases selected:</u> Multiple databases	
Document View « Back to Results	Document 1 of 5 Next >
Print Email Mark Document	Publisher Information
None of your e-business	Other available formats:
Aviel D Rubin. Web Techniques. San Francisco: Apr 20 Vol.5, Iss. 4; pg. 55, 4 pgs	<u>Abstract</u> □
» Jump to full text	Full Text
» Translate document from: Select language	The Full Text - PDE

Subjects:

Computer security, Electronic

commerce, Information, Computer privacy

Locations:

United States, US

Author(s):

Aviel D Rubin

Document types:

Feature

» More Like This - Find similar documents

Publication title:

Web Techniques. San Francisco: Apr

2000. Vol. 5, Iss. 4; pg. 55, 4 pgs

Source type:

Periodical

ISSN:

1086556X

ProQuest document ID: 50708866

Text Word Count

3189

Document URL:

http://proquest.umi.com/pqdweb?did=50708866&sid=5

&Fmt=4&clientId=19649&RQT=309&VName=PQD

Abstract (Document Summary)

Online credit card information is potentially available to many people who can access it, either legitimately or otherwise, from remote locations, over networks that may or may be secure. Surprisingly, many people still believe that their e-mail is somewhat private. Encryption programs are a viable solution provided that they are installed by both parties. Direct compromise of privacy occurs when a company actively collects personal information, then indexes, categorizes, and shares it with people. Public humiliation is a great way to teach companies that blatantly invade people's privacy a lesson and get them to stop. Indirect methods involve compromises of privacy that are a result of a security incident. One way to ward off the indirect attacks described is to protect the data on a machine. People can protect their credit card numbers by refusing to enter credit card information into a Web form. There are tools for protecting browsing activity from Web servers and other users.

Full Text (3189 words)

Copyright Miller Freeman Inc. Apr 2000

The advent of the Internet and Its wide-scale adoption have irrevocably changed the way we interact with each other, the way we transact

business, and the way personal information is gathered and maintained about us. While progress is always exciting and provides a multitude of opportunities, there's also a dark side. Never before has there been such an opportunity for aggregating and cross-referencing personal information about people.

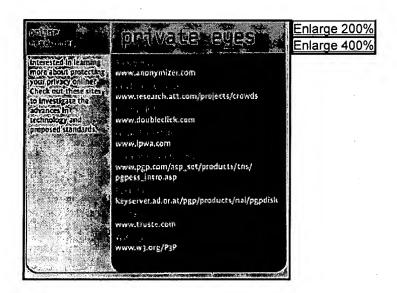
Before the Internet, when you physically browsed a bookstore, the store might keep track of which books you bought and which you returned. Now, you're more likely to shop online at amazon.com or barnesandnoble.com. These stores can keep track of which books you bought, which books you browsed, from which books you read a sample chapter, which books you recommended, and which books you avoided. They even know what search terms you use to locate books. Similarly, the online grocery can elicit much more information from its customers than the physical supermarket of the past, and so can the pharmacy and the online video store. In short, when you interact online, a record is created of every single action you take. But this is only half of the picture. More and more organizations are going online. So the sanctity of the sensitive records they keep is subject to the quality of their computer security. It doesn't do any good for an organization to maintain a respectable privacy policy if its machines are wide open to attack. Unfortunately, security on the web is dismal at best.

As if accidental disclosure of personal information weren't enough, there's also a huge incentive for organizations to share the information they collect. Targeted advertising is an enormous and growing market. There are companies whose sole purpose is to collect as much personal information as possible about people, cross-reference it, and sell the results to other companies.

What Is at Risk?

The privacy of many types of personal data can be compromised. Perhaps the most personal and valuable thing that you have is your identity. This includes your name, your social **security** number, and the profile kept by the organizations with which you interact. Disclosure of this data can result in theft of your identity. While this concept has been discussed hypothetically in books and movies (for example, The Net), the threat is quite real.

Other information is at risk as well. Your credit card numbers are hardly private. You give them out to the waiter at the restaurant, the telephone agent of the airline, your grocery store, and all of your online retailers. On the other hand, in all cases, these parties have an incentive to keep the credit card number to themselves. Unfortunately, it's very easy for this information to end up in the wrong hands. Online credit card information is potentially available to many people who can access it, either legitimately or otherwise, from remote locations, over networks that may or may not be secure. Credit card information sits in databases on computers that are often much more vulnerable than their owners realize. Recent high-profile extortion cases, where hackers demanded ransom from banks by threatening to release their customers' credit card information, demonstrate how personal information can be compromised.



You may believe that your browsing habits-namely, which Web sites you visit, and what you do there-are private. You may believe that data on your computer is private. You may believe that the email you send is private. But all of these things are at risk on the Internet. We'll look at ways in which such private information is compromised and countermeasures that are available.

Email Privacy

With the possible exception of Web browsing, email is probably the most commonly used network application. If you're reading this magazine, you don't need to be convinced that email is as ubiquitous as the telephone. For many people, it may be even more so. if you ask most people, they'll probably tell you that their company has the ability, and in some cases the right, to read all email messages. What they probably don't realize is the extent to which email monitoring technology has advanced. Corporations can use tools to regularly monitor email for personal messages, messages that could indicate inappropriate behavior, and other information that the company could find to be of use. These tools allow sophisticated keyword searches, as well as complicated heuristics for automatically sifting through email to find whatever the authorities decide they need to find.

Surprisingly, many people still believe that their email is somewhat private. Messages are sent with personal information, ranging from gossip to performance review data, that should never be seen by other employees. Most employees don't think about the fact that not only do the system administrators have access to all email on the system, but, in many networks, so does just about everyone else.

The most common type of network in the office environment is the Ethernet. It's trivial to tap into the Internet and monitor every packet that travels on the LAN. Available programs can take the raw IP packets and reconstruct TCP streams. From there, application-level data can be assembled, and email can be monitored. Programs called sniffers can be configured easily to capture every email message that travels on an Ethernet.

Fortunately, you can take steps to counter the threats to email privacy. Unfortunately, none of the steps is that simple. The problem is that email is shared between at least two people. It doesn't make sense to protect email messages cryptographically if the recipient of your message doesn't have the ability to decrypt. Unless you two share a key, and the software to

perform encryption and decryption, your message is going to be very private-so private that even your intended recipient won't be able to read it.

Encryption programs are a viable solution provided that they're installed by both parties. One good example is PGP, which uses the RSA public key cryptosystem to provide digital signature and encryption capabilities for email. To use the system, you must first generate a key pair, using the software; the other party must do the same. Next the public keys are exported into a file, where they are manually exchanged by both parties. "Manually" means that you shouldn't use an untrusted network to exchange the keys. The best way to exchange them is in person. If this isn't practical, you can exchange the public keys over an insecure network, but after doing so, verify the built-in fingerprint of each key with each other by phone. The fingerprint is simply the MD5 hash of the bits of the public key, in hex format. MD5 is a cryptographic checksum function that is used to detect data tampering.

Once you have each other's public keys, you can use the PGP software to encrypt and/or sign a message. On the other end, the software can bi used to decrypt and/or verify signatures, depending on the option chosen by the sender.

PGP is just one option. Another solution that also uses the RSA cryptosystem has been standardized by the Internet Engineering Task Force (IETF). This is called S/MIME, and it's convenient in that it's nicely integrated into the commonly used Microsoft Outlook and Netscape Messenger mail programs. To use S/MIME, you must first generate a key pair and obtain a certificate for your public key. The good thing about a standard is that the software for using it is available in many different browsers. The browsers also contain root public keys that can be used to verify a certificate, regardless of which certifying authority signs it.

The main difference between PGP and S/MIME is that in the former, users are responsible for distributing their public keys to one another. on the other hand, in S/MIME, certifying authorities act as introducers for people, and as long as everyone shares a common set of root public keys, the authentication system works. PGP is more appropriate for a diverse set of people with no common authority. S/MIME is geared toward organizations with tight hierarchical structures.

Other systems, such as Lotus Notes, operate similarly to S/MIME. However, users should be careful when using a mail system in which the cryptography is administered by their own organization. It's possible, even likely, that management has a way of bypassing the cryptography to get at the messages. This can be accomplished by making extra copies of private keys, or by using functions with back doors.

How Is Personal Privacy Compromised?

Whether you're concerned with email privacy, confidentiality of credit card numbers, or protecting your browsing history, you need to understand how the risks manifest themselves directly and indirectly. Direct methods are those that result in loss of privacy due to action on the part of the parties with whom you interact and their partners. Indirect methods are ways in which privacy can be lost despite the best behavior of the parties involved.

The Direct Approach

Direct compromise of privacy occurs when a company actively collects your

personal information, then indexes, categorizes, and shares it with people. This has been going on for a long time. Personal information about you is so valuable that many organizations offer free or discount services in exchange for this information. The Web is an ideal platform for companies that operate in this space. Take DoubleClick, for example. it stores ads on Web pages and then uses cookies to correlate requests across Web sessions.

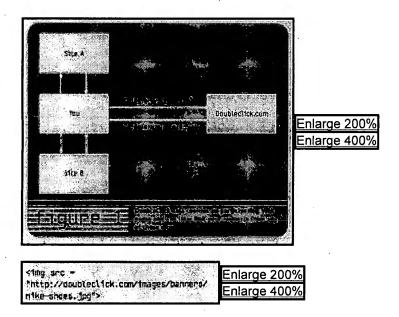


Figure 1 illustrates how this works. You, the innocent user, visit site A. in doing so, you download a Web page from the site. The Web page that you download has, among other things, an image that resides on DoubleClick.com's Web site. For example, the following HTML segment

Set-cookte: vatue=3452385rkA35rtk]874582 Enlarge 200%

causes your browser to fetch the image nike-shoes.jpg from the DoubleClick site. When DoubleClick serves the page, it also includes a cookie in the response. This is just a **header** of the format:

Upon receiving this, your browser displays the image, and stores the cookie value in the cookie file on your local disk.

At some point in the future, you visit site B. This site is also a customer of DoubleClick's. When your browser loads site B's Web page, it contains an embedded image from DoubleClick. When the page loads, your browser requests the image. However, since your cookie file contains a cookie that was set by DoubleClick, the cookie is returned as well. Thus, DoubleClick can correlate that the request from site A and the request from site B came from the same browser. As a result, DoubleClick can build a whole list of sites that a particular user visits. This is due to a special HTTP header called the Referer field. The Referer head relays to DoubleClick the URL of the page that the user requested from site A or B. Consider the implications of this. Many search engine requests include the search terms in the URL. In fact, when search values are sent to a CGI script on a server using the GET method, the fields filled in by the user are all included in the URL. This gives DoubleClick all sorts of valuable information, which can be mined by powerful algorithms and offered for sale to the highest bidder. For example, a company that determines whether to approve someone for health

insurance might be interested in this information.

Another example of surreptitious use of the direct method involves RealNetworks. This company provides streaming video and audio through Web browsers. After it was in business for some time, an independent security expert, Richard Smith, discovered that the RealNetworks software collected information from peoples' browserswithout their permission-about the content that users requested. This invasion of privacy was halted when stories in the media proved a big slap in the face to the company. Public humiliation is a great way to teach companies that blatantly invade your privacy a lesson, and more importantly, to get them to stop.

Covert Operations

Indirect methods involve compromises of privacy that are a result of a security incident. For example, a malicious worm can propagate quickly across the Internet. Worms such as Melissa, Explorer.zip, and friends may have been infuriating, but as far as we know, they weren't attacks against privacy. It isn't too difficult to envision a malicious version of a worm that secretly searches a computer for "interesting" looking content and covertly discloses it to a remote location. While we haven't seen such a thing on a broad scale, these types of malicious programs could be a very real danger in the near future.

OK, enough gloom and doom. what can you do to protect yourself? We'll examine some defense mechanisms you can use to try to win back some of your privacy.

Protect Data an Your Machine

One way to ward off the indirect attacks described is to protect the data on your machine. You want to protect the information for the worstcase scenario-such as when a perpetrator runs arbitrary programs, like an Internet worm/virus, right on your computer. The only way to properly achieve this is with cryptography. There are programs available for encrypting the information on your hard drive. A great one is PGPDisk, which lets you mount an encrypted drive. The data is encrypted with a key derived from a user-supplied passphrase. When the drive is mounted, all of the data is decrypted as it is read from disk and encrypted as it is written to disk. When the drive isn't mounted, the data is available only in encrypted form, and the key is removed from memory. The privacy of your information is preserved because an intruder has no way of reading your files.

Protect Credit Card Information

Credit card information is the most widely used personal data that users readily provide to Web servers. Recent cases of extortion demonstrate that lists of credit card numbers are very valuable to merchants, advertisers, the owners of the cards, and, probably most of all, the credit card companies. The inherent insecurity of the Internet makes it relatively easy for a determined attacker to steal credit card numbers. What can you do to protect your credit card numbers? There are several approaches.

The first thing you can do is refuse to enter your credit card information into a Web form. While this may seem drastic at first, you can be assured that most Web merchants offer a mechanism for you to provide the payment information by telephone or mail. Keep in mind that the only insecure link that you avoid by doing this is the connection from your computer to the server. All of the back-end processes on the server will still contain your

information as soon as you speak it into the phone.

If you're a bit more intrepid, enter your credit card number into a Web form. But first, check to make sure that you're communicating over a secure channel that uses SSL. Look for a closed lock on the browser and make sure that the URL is of the form https:// as opposed to http://. Once you know that you're using SSL, you can be pretty confident that the session is encrypted. To avoid poor encryption, you should examine your security preferences in the browser and turn off SSL version 2, as well as all of the ciphersuites that use 56-bit encryption or less. Next, it's important to verify that your machine is speaking with the right server and not an imposter that's trying to get at your personal information. To verify this, before you enter your credit card number, check that the certificate is signed by an authority that you know and trust, and that the certified entity in the certificate is the one that you want to provide with your credit card number. Finally, pay attention to any warnings that the browser pops up at you.

Although your personal liability is probably limited to \$50 in case of credit card fraud, remember that other kinds of misuse of your private information may go unnoticed for a while. It's important to be very careful when using "secure" Web sites. Keep in mind that even if the channel from you to the merchant is secure, you have almost no way of knowing how carefully the merchant's computers are monitored.

One final word on protecting credit card information. Check whether the sites with which you shop offer a privacy policy. Read the policy and make sure that you're willing to accept the terms listed. There are privacy sea] programs, such as TrustE, that audit sites to make sure they're complying with their stated policies. If you believe that they're doing a good job, it may give you reason to have more confidence in a privacy policy. Finally, a more long-term effort is underway at the W3C consortium. A project called P3P is designed to allow users and Web sites to automatically negotiate a level of privacy protection for information. (For more information, see the article "Privacy Critics," Web Techniques, September 1999.)

Safeguard Your History

There are several ways that your browser keeps track of your online browsing habits. The browser history shows past URLs you've requested. In Netscape Navigator, the prefs is file keeps a list of URLs you entered into the location window. Also, the local cache keeps copies of things you browse. if you share your computer with others, it's a good idea to clear all of these if you care to protect your privacy regarding the sites that you browse.

There are also tools for protecting your browsing activity from Web servers and other users. AT&T Labs Crowds is a system for browsing the Web in such a way that others can't tell that it's you browsing. To use the system, you join a crowd of other users by running a **proxy** on your machine. From then on, your actions are indistinguishable from those of others in the crowd. Another option called Anonymizer is a commercial system that hides your location from end servers. However, the administrators of the Anonymizer themselves can collect the information. Another system, ProxyMate, can be used for Web sites that require subscription services. ProxyMate provides you with pseudonyms and remembers passwords for you. The names and passwords are generated randomly, so that Web sites can't link your activity to a real-world person based on the information you enter.

Summary

In our world of ever-advancing technology, privacy has become a casualty. As more of our daily activities, interests, and relationships move onto computers, the potential for misuse of this information increases. Web server logs do not forget, and the technology for correlating information and building up data shadows about people is constantly improving. The technologies are in their infancy, and the ability to compromise privacy is growing faster than our ability to protect ourselves. Without a focused research effort on privacy technologies, and perhaps legislation, the future may be a much darker place. :0-C

[Author Affiliation]

More Like This - Find similar documents Subjects: Computer security Electronic commerce ■ Information ■ Computer privacy Locations: Author(s): Aviel D Rubin Document types: Feature Language: English Publication title: Web Techniques Clear Search Mark Document Print **Email** Publisher Information Back to Top « Back to Results Document 1 of 5 Next >

Copyright © 2007 ProQuest-CSA LLC. All rights reserved.



Pro Quest

Return to the USPTO NPL Page | Help

В	asic		Advanced	Topics	Publications	My R]	English
Data	bases	sel	ected: Multiple	databases	•				
Re	sult	:S -	powered by P	roQuest [®] Sm	nart Search				
	Sugg	este	ed Topics Ab	<u>out</u>	•			•	< Pri
	Secu	rities	regulations Al	ND Proxies					
			·						
									•
		•	•			-			
			ts found for: F HTTP and hea		(2000) and (pro	oxy and (sec	ure or	» Refine	Search Set Up Alert 🖂
All	sourc	ęs	Scholar Scholar	ly Journals	<u>Magazines</u>	Trade Publi	<u>cations</u>	<u>Newspapers</u>	<u> </u>
	Mark	k all	🗁 0 marke	ed items: Em	ail / Cite / Expo	ort [Show	only full text	Sort results by: Most n
	1.	Tu	rn a Solaris b	oox into a p	acket-filtering	firewall	lee den	1 (6 pages)	
		БО	ris Loza. Insid		<u>_</u>	2000. Vol. 6, <u>ull Text - PD</u> F	•	. i (o pages)	et
		 Co		and the second of the second s	ronic informa	er weldtig sjoer oorganise tijder en were tij en trijder ve troes progress tijde tijde	seeme to the constant of the c	g gillion, den å likedikkilliskilliskillingsreperiskilling av av av dillion.	tining the state of the state o
	۷.						User Co	mputing. Apr-	lun 2000. Vol. 12, Iss. 2; p.
	·		B <u>Text+C</u>	<u> Braphics</u>	₺ <u>Fu</u>	ıll Text - PDF	<u> </u>	Abstrac	<u>ct</u>
	3.	No Av	ne of your e- iel D Rubin. W	<u>business</u> /eb Technic	lues. Apr 2000). Vol. 5, Iss.	4; p. 55	(4 pages)	comming and the state and states about abstract and the semigroups and semigroups.
		ı	Text+C	<u> Braphics</u>	<u> </u>	<u>ıll Text - PDF</u>	<u>:</u>	Abstrac	<u>ct</u>
	4.		intaining sec f Forristal. Ne			asset: Mar 20	, 2000. \	/ol. 11, lss. 5; p	o. 93 (4 pages)
			Text+C	<u> Braphics</u>	₺ <u>F</u> u	ıll Text - PDF	<u>.</u>	<u>Citatior</u>	1
固	5.		Framework a iro A. Gutierre rk: Mar 2000.	z, Donald P.		for Multimed Radhakrishn	dia Netw a Pillai. J	ork Managemolournal of Netv	ent work and Systems Manage
			Article	<u>image - PDf</u>	<u>.</u>			<u>Abstract</u>	<u>ct</u>
	6.	<u>Ca</u> Joi	ching in with nathan Angel.	content de Network M	<u>livery</u> agazine. Mar 2	2000. Vol. 15	i, Iss. 3; p	o. 92 (5 pages)	enter en enternacione de la companya
			Text+C	<u> Braphics</u>	D FL	ıll Text - PDF		Abstrac	<u>ct</u>
Ø	7.		rnkey solutio Agranat, Dar			ngineering T	imes. Ma	anhasset: Feb	28, 2000. ; p. 106 (4 pages)
			Text+C	<u>Graphics</u>	ሺ <u>F</u> ∟	<u>ıll Text - PDF</u>	Ξ.	<u>Citatior</u>	<u>1</u>
	8.		ookie cutting d Schroeder.			am: Feb 21,	2000. Vo	ol. 17, Iss. 8; p.	49 (1 page)
			Text+C	<u>Graphics</u>	₺ <u>F</u> L	<u>ıll Text - PDF</u>	<u> </u>	Abstrac	<u>ct</u>

Б	9.	IO-Lite; a unified I/O buffering Vivek S. Pai, Peter Druschel, W 18, Iss. 1, p. 37	<u>g and caching system</u> /illy Zwaenepoel. ACM Transact	ions on Computer Systems. New York: F
		Link to full text	•	Abstract
	10.	My agent will call your agent Martin L Griss. Software Devel	opment. San Francisco: Feb 20	00. Vol. 8, Iss. 2; p. 43 (4 pages)
				Abstract
	11.	IP challenges QoS requireme Azhar Sayeed. Electronic Eng	nts ineering Times. Manhasset: Jar	n 24, 2000. ; p. 110 (4 pages)
		<u>Text+Graphics</u>	D Full Text - PDF	[™] <u>Citation</u>
	12.	Web traffic complexitites cha Loring Wirbel. Electronic Engli	nge LANscape neering Times. Manhasset: Jan	24, 2000. ; p. 93 (2 pages)
		Text+Graphics	D Full Text - PDF	in Citation
Ø	13.		MMON PROTECTION SCHEMES Fribune. Chicago Tribune. Chica	
		Full text		Abstract
	14.		to address critical security iss rity Journal. San Francisco: Win	ues ter 2000. Vol. 16, Iss. 1; p. 15 (24 pages)
				Abstract
	15.		for optimizing emulator perforr tems Journal. Armonk: 2000. Vo	
		™ Text+Graphics	D Full Text - PDF	[™] Abstract
	16.	A model for library support of Jean L.Cooper. Interlending &	f distance education in the USA Document Supply. Bradford: 20	<u>A</u> 000. Vol. 28, Iss. 3; p. 123
		Full text		Abstract
W	17.		upply: a review of recent literat cument Supply. Bradford: 2000	
		Full text		Abstract Abstract
	18.	Rawn Shah. JavaWorld. San F	to take over; WAP addresses rancisco: Jan 1, 2000.; p. 1	the shortcomings of other protocols
		Full text		Abstract
	19.	Bobby: An online tool for Wel Jessica Chaiken, David Johnso	<u>b site designers</u> n. Library Hi Tech News. Bradfo	ord: 2000. Vol. 17, lss. 2; p. 28 (4 pages)
		Text+Graphics	D Full Text - PDF	Abstract
	20.		manage network resources, it of Systems Journal. San Francisco	can manage your data too b: Jan 2000. Vol. 15, Iss. 1; p. 51 (14 page: Abstract
		CONTRACTOR OF THE PROPERTY OF THE SECTION OF THE SE	ия от та этой го этоми и переду при	— AUSTRALIA
	21.	On proxy agents, mobility, an Anupam Joshi. Mobile Network	id web access ks and Applications. Amsterdan	n: 2000. Vol. 5, Iss. 4; p. 233
		Article image - PDF		[™] <u>Abstract</u>

	22.			ents and services ik Lamming. Mobile Networks and Applica
		Article image - PDF		Abstract Abstract
	23.	Optimizing Internet bandwing Ralph Barker. Performance		000. Vol. 18, Iss. 2; p. 29 (5 pages)
		Text+Graphics	D Full Text - PDF	[⊕] <u>Abstract</u>
	24.	Sites Strive To Hold On To Christine Zimmerman. Intern	Buyers etWeek. Manhasset: Nov 29, 1999	9. ; p. PG.1
		Full text		[⊞] Abstract
	25.			on seph, Randy H. Katz. IT Professional Mag
٠		D Full Text - PDF		Abstract Abstract
	26.		ting and using networked inform Tech News. Bradford: Nov 1999.	
		Full text	D Full Text - PDF	[⊞] Abstract
	27.		s: Expectations, ideals and realing ecurity Journal. San Francisco: F	ties Fall 1999. Vol. 15, Iss. 4; p. 25 (21 pages)
	28.	The QoS quagmire Tom Stenson. Network Worl	d. Framingham: Sep 6, 1999. Vol.	16, Iss. 36; p. 53 (3 pages)
		Text+Graphics	D Full Text - PDF	[™] Abstract
	29.	WAP: New Web whopper Joshua Piven. Computer Tec	chnology Review. Los Angeles: S	Sep 1999. Vol. 19, Iss. 9; p. 1 (3 pages)
			D Full Text - PDF	[™] Abstract
	30.	Video servers revisited Jonathan Angel. Network Ma	gazine. Sep 1999. Vol. 14, Iss. 9;	p. 56 (5 pages)
		Text+Graphics	D Full Text - PDF	[™] Abstract
1-30	of 15	от при на при	 - Симомутичной станстватумы и постоя постоя	< First < Previous 1 2
		()	this search? Set Up Alert	Results p
			If not, <u>refine your search</u> below o	
		ested Topics About		< Pri
	Secu	rities regulations AND Proxies	•	

△Tools: Search Tips Browse Topics 2 Recent Searches △

PDN(<04/10/2000) a	and (proxy and (secure or security) and HTTP and heade Search Clear
Database: Mult	iple databases Select multiple databases ates
Limit results to: 🖫 F	ull text documents only
□ s	cholarly journals, including peer-reviewed 🎓 About
More Search Options Publication title:	^ Hide options Browse publications About
Author:	About
Look for terms in:	Citation and document text About
Document type:	Any document type
Publication type:	All publication types
Exclude from results:	Book Reviews
	☐ Dissertations
	Newspapers
Sort results by:	Most recent first

Copyright © 2007 ProQuest-CSA LLC. All rights reserved.

